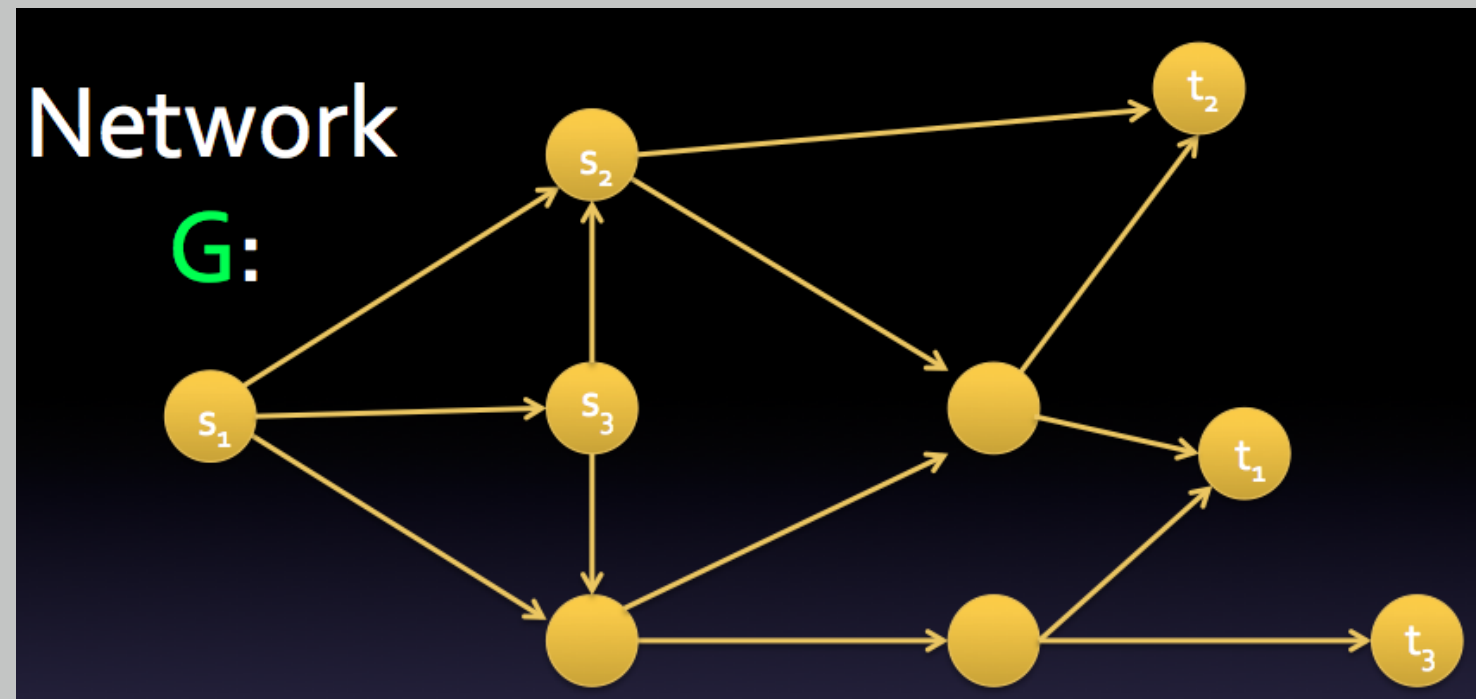


Approximately Optimal Flow via Truthful Mediators



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Socially Optimal Flow in Routing Games



- Losses on each edge $\ell_e(y_e)$.
- Player i routes one unit of flow from s_i to t_i .

Main Assumption - Large Games



Large Game

Any player has a small $o(1)$ impact on the costs of others as $n \rightarrow \infty$.

Classical Approach - Impose Tolls



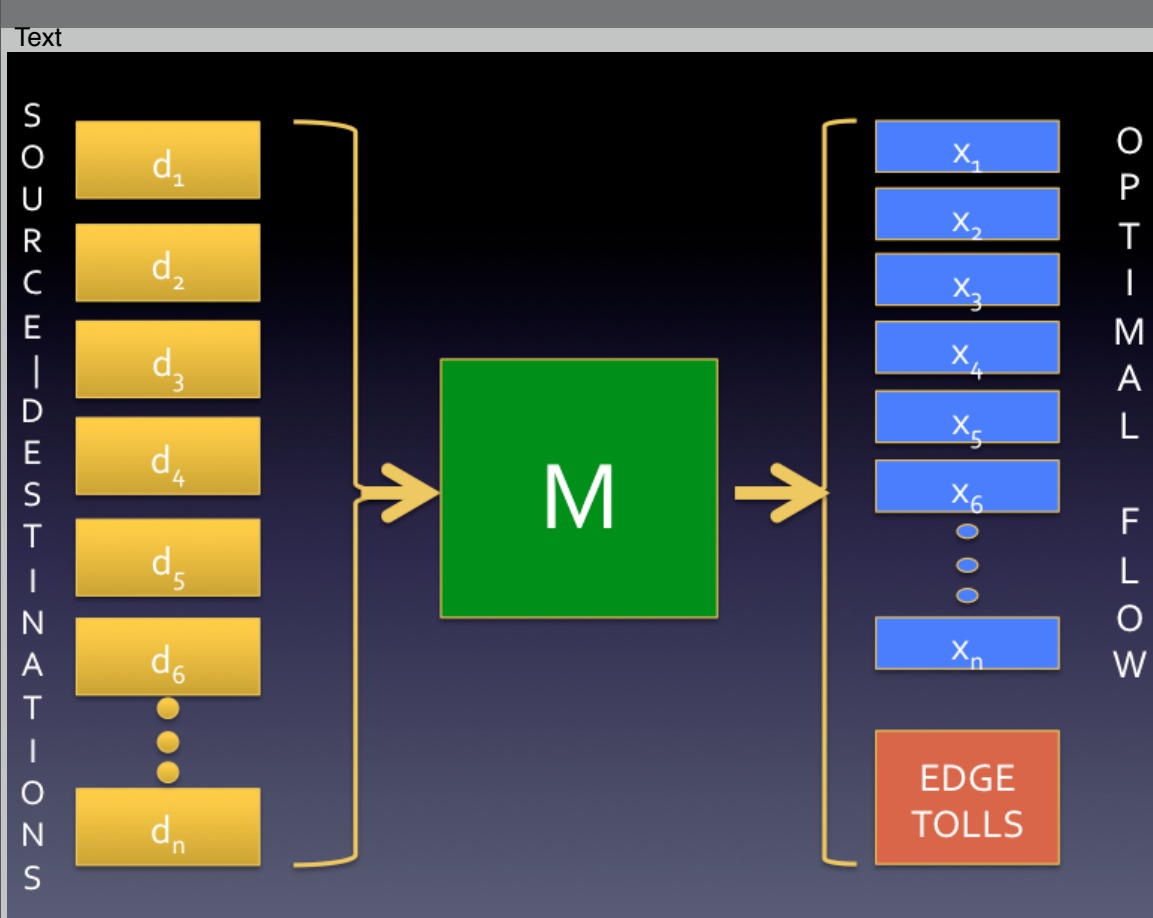
- Want a socially optimal flow
- A Mediator M may enforce tolls on each edge so selfish players route optimally.

Enter Mechanism Design

- The tolls the mediator computes depends on the demands.
- Want players to truthfully report their demands so mediator can compute the correct tolls



Introduce a Mediator that can Enforce Tolls



Weakly Mediated Game

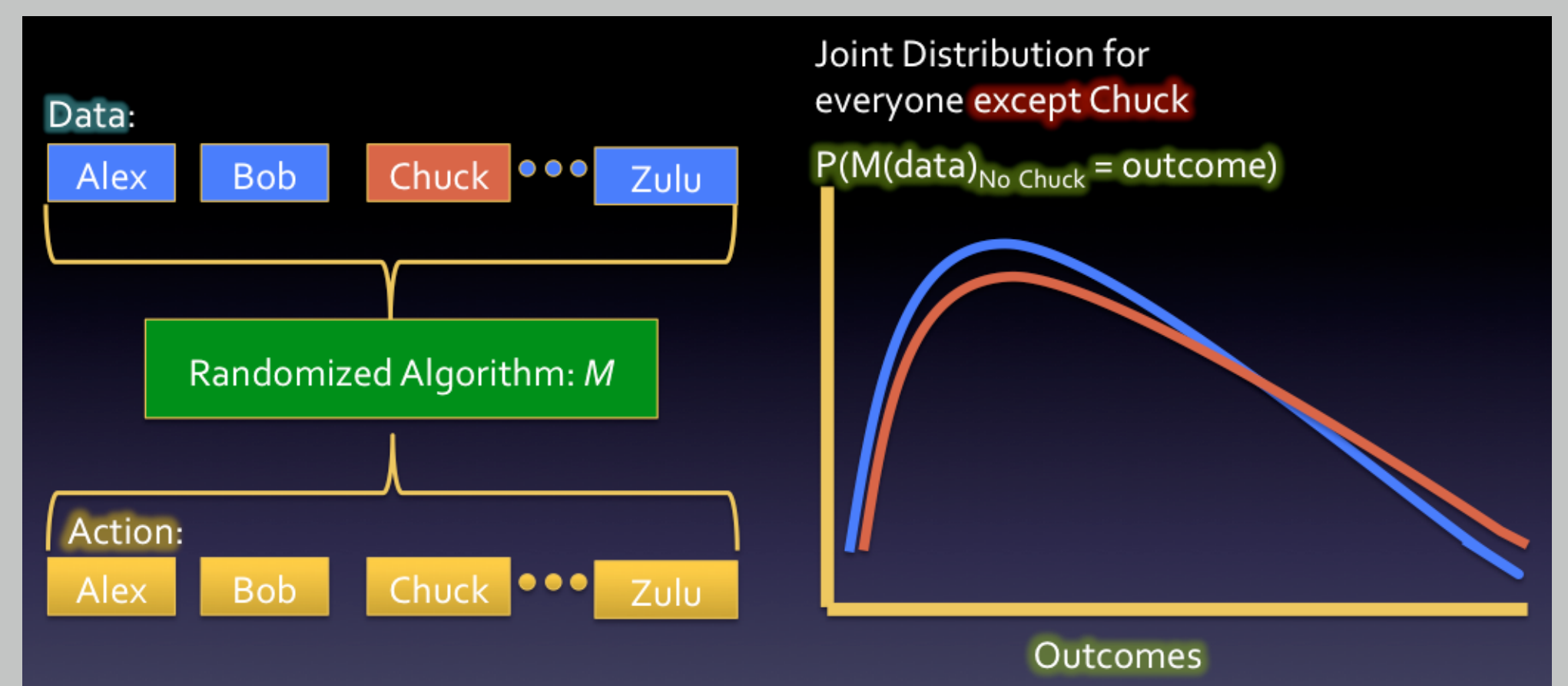
- May bypass M .
- May misreport to M .
- May not follow M 's suggested route.

Main Result

- We develop a mediator such that for $n \rightarrow \infty$
 - Reporting truthfully and following the suggested action of M is an (asymptotic) **ex - post Nash equilibrium**
 - The resulting flow has cost $(1 + o(1))OPT$

Useful Tool - Joint Differential Privacy

Joint Differential Privacy [KPRU'14]

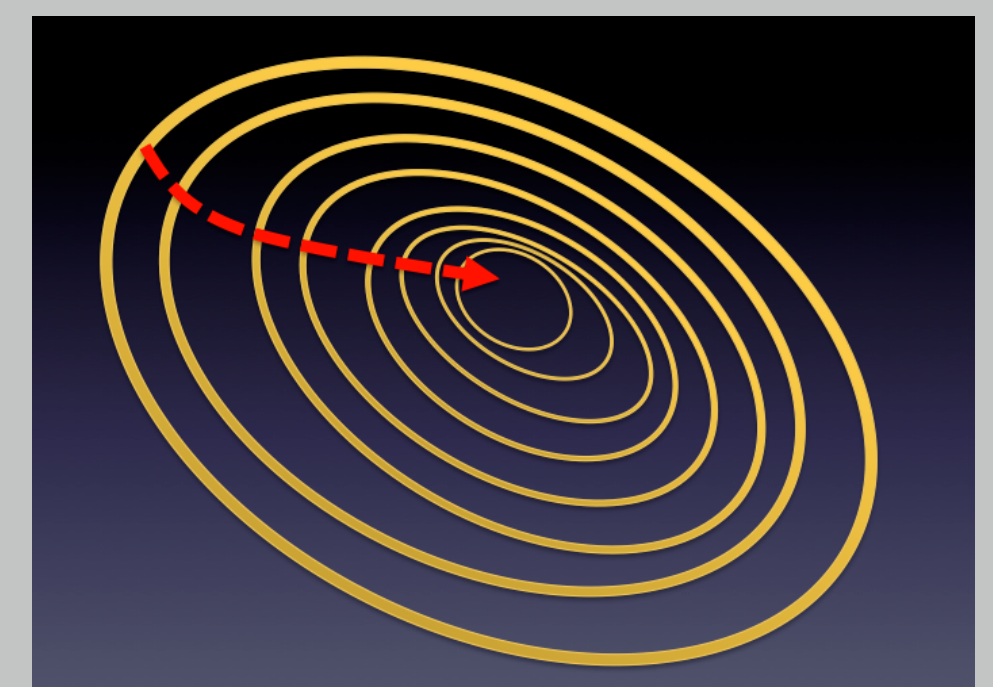


- Controls the impact a single player has on the outcome for other players.
- No real "privacy" concerns here, but still useful!

Novel Technique - Private Gradient Descent

We need to solve the convex program in a way that is joint differentially private in data s .

$$\begin{aligned} \min \quad & \text{Total Cost of } x \\ \text{s.t.} \quad & x \in \mathcal{F}(s) \end{aligned}$$



Conclusion

We construct a weak mediator M in a **Large** Routing Game that adds **tolls** to induce optimal flow, with **unknown** player demands.

Open Problem

Can we make weak mediators in **general** large games?