CSci 5271 Introduction to Computer Security Day 25: Electronic cash and Bitcoin

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Outline

Cryptography for voting

Previous e-cash and techniques

Announcements intermission

Bitcoin design

Bitcoin experience



Commitment to values

- Two phases: commit, later open
 Another analogy to a use of envelopes
- Binding property: can only commit to a single value
- Hiding property: value not revealed until opened
- Trivia: either binding or hiding, but not both, can be perfect
 - Information-theoretic, like a one-time pad

Randomized auditing

- How can I prove what's in the envelope without opening it?
- n envelopes, you pick one and open the rest
 - \blacksquare Chance 1/n of successful cheating
- Better protection with repetition

Election mix-nets

- Independent election authorities similar to remailers or Tor nodes
- Onion-encrypt ballot, each authority shuffles and decrypts
- Extra twist: prove no ballots added or removed, without revealing permutation
 Instance of "zero-knowledge proof"
- Privacy preserved as long as at least one authority is honest

Pattern voting attack

- Widely applicable against techniques that reveal whole (anonymized) ballots)
- Even a single race, if choices have enough entropy

3-choice IRV with 35 candidates: 15 bits

- Buyer says: vote first for Bob, then 2nd and 3rd for Kenny and Xavier
 - Chosen so ballot is unique

Fun tricks with paper: visual crypto

- Want to avoid trusted client, but voters can't do computations by hand
- Analogues to crypto primitives using physical objects
- One-time pad using transparencies:



Scantegrity II Designed as end-to-end add-on to optical scan system Fun with paper 2: invisible ink Single trusted shuffle Checked by random audits of commitments Version used in a DC-suburb municipal election

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Wide adoption among consumers, little

consumer fraud liability ■ Restrictive merchant procedures

💼 PayPal

- Easier to accept payments
- Centrally managed to deal with fraud

Ideal: electronic cash

- Direct transactions without third party
- 🖲 No transaction fees
- Potentially anonymous
- Non-revocable: buyer bears fraud risk



Blinded signatures Sign something without knowing its value Often used together with randomized auditing For RSA, multiple message by r^e, r random Allows a bank to "mint" coins that can still be anonymous





Hash trees and timestamp services

- Merkle tree: parent node includes hash of children
- **Output** Good hash function \rightarrow root determines whole tree
- Can prove value of leaf with log-sized evidence
- Application: document timestamping (commitment) service

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HW2 due Sunday

- Non-early due date: 11:55pm this Sunday
- Q5 performance/load issues
 - Avoid by not doing Q5 at the last minute, testing on yourself



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Bitcoin addresses Address is basically a public/private signing key pair Randomized naming, collision unlikely At any moment, balance is a perhaps fraction number of bitcoins (BTC)

Anyone one can send to an address, private key needed to spend



Bitcoin network

- Use peer-to-peer network to distribute transaction log
- Roughly similar to BitTorrent, etc. for old data
- Once a client is in sync, only updates need to be sent
- New transactions sent broadcast

Consistency and double-spending If all clients always saw the same log, double-spending would be impossible But how to ensure consistency, if multiple clients update at once? Symmetric situation: me and "me" in Australia both try to spend the same \$100 at the same time

Bitcoin blocks

- Group ~10 minutes of latest transactions into one "block"
- Use a proof of work so creating a block is very hard
- All clients race, winning block propagates

Bitcoin blockchains

- Each block contains a pointer to the previous one
- Clients prefer the longest chain they know
- E.g., inconsistency usually resolved by next block

Regulating difficulty

- Difficulty of the proof-of-work is adjusted to target the 10 minute block frequency
- Recomputed over two-week (2016 block) average
- Network adjusts to amount of computing power available

Bitcoin mining

- Where do bitcoins come from originally?
- Fixed number created per block, assigned by the client that made it
- Incentive to compete in the block generation race
- Called mining by analogy with gold

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- 🖲 Block chain 271,000 blocks, about 14GB
- 12M BTC minted (many presumed lost)
- Theoretical value at market exchange rate > \$1 billion
- Millions of addresses, probably many fewer users
- Mining power: 5 petahash/sec







Bitcoin mining trends

- Exponentially increasing rates
- \bigcirc CPU \rightarrow GPU \rightarrow FPGA \rightarrow ASIC
- Specialized hardware eclipsing general purpose
 - Including malware and botnets
- Recent price trends suggest continuing investment



