CSCI 2021: Binary Floating Point Numbers

Chris Kauffman

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Logistics

Reading Bryant/O'Hallaron

- Ch 2.4-5: Floats, Wed/Fri
- 2021 Quick Guide to GDB
- Next week: Ch 3.1-7: Assembly Intro

Goals this Week

- Discuss Bitwise ops from Integer Rep Slides
- Floating Point layout
- gdb introduction

Assignments

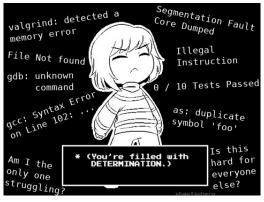
- Canvas Feedback Survey
 - Anonymous: be honest!
 - Worth 1 EP
 - Due Wed 15-Feb
 - ▶ 67% response rate so far
- ▶ HW04: Due Wed 11:59pm
- Lab05: Bit operations
- ► HW05: Bits, Floats, GDB

P2 Released

- Bit shifting and Debugger Usage
- Due date pushed back to Mon 27-Feb

P1 'sanity' submission Problems See Piazza announcement here: https://piazza.com/class/lcsjsmrfvdb1k4/post/201

Don't Give Up, Stay Determined!



- If Project 1 / Exam 1 went awesome, count yourself lucky
- If things did not go well, Don't Give Up
- Spend some time contemplating why things didn't go well, talk to course staff about it, learn from any mistakes
- There is a LOT of semester left and plenty of time to recover from a bad start

Parts of a Fractional Number

The meaning of the "decimal point" is as follows:

$$123.406_{10} = 1 \times 10^2 + 2 \times 10^1 + 3 \times 10^0 + 123 = 100 + 20 + 3$$
$$4 \times 10^{-1} + 0 \times 10^{-2} + 6 \times 10^{-3} \quad 0.406 = \frac{4}{10} + \frac{6}{1000}$$
$$= 123.406_{10}$$

Changing to base 2 induces a "binary point" with similar meaning:

$$110.101_{2} = 1 \times 2^{2} + 1 \times 2^{1} + 0 \times 2^{0} + \qquad 6 = 4 + 2$$
$$1 \times 2^{-1} + 0 \times 2^{-2} + 1 \times 2^{-3} \qquad 0.625 = \frac{1}{2} + \frac{1}{8}$$
$$= 6.625_{10}$$

One could represent fractional numbers with a fixed point e.g.

- 32 bit fractional number with
- 10 bits left of Binary Point (integer part)
- 22 bits right of Binary Point (fractional part)

BUT most applications require a more flexible scheme

Scientific Notation for Numbers

"Scientific" or "Engineering" notation for numbers with a fractional part is

| Standard | Scientific | <pre>printf("%.4e",x);</pre> |
|----------|--------------------------|------------------------------|
| 123.456 | 1.23456×10^{2} | 1.2346e+02 |
| 50.01 | $5.001 	imes 10^1$ | 5.0010e+01 |
| 3.14159 | $3.14159 	imes 10^{0}$ | 3.1416e+00 |
| 0.54321 | $5.4321	imes10^{-1}$ | 5.4321e-01 |
| 0.00789 | $7.89	imes10^{-3}$ | 7.8900e-03 |

- Always includes one non-zero digit left of decimal place
- Has some significant digits after the decimal place
- Multiplies by a power of 10 to get actual number

Binary Floating Point Layout Uses Scientific Convention

- Some bits for integer/fractional part
- Some bits for exponent part
- All in base 2: 1's and 0's, powers of 2

Conversion Example

Below steps convert a decimal number to a fractional binary number equivalent then adjusts to scientific representation.

```
float fl = -248.75;
```

```
7 6 5 4 3 2 1 0 -1 -2
-248.75 = -(128+64+32+16+8+0+0+0) \cdot (1/2+1/4)
        = -11111000.11 * 2^{0}
            76543210 12
        = -1111100.011 *2^{1}
            6543210 123
        = -111110.0011 *2^{2}
            543210 1234
              MANTISSA
                           EXPONENT
        = -1.111100011 * 2^{7}
            0 123456789
Mantissa \equiv Significand \equiv Fractional Part
```

7

Principle and Practice of Binary Floating Point Numbers

- In early computing, computer manufacturers used similar principles for floating point numbers but varied specifics
- Example of Early float data/hardware
 - Univac: 36 bits, 1-bit sign, 8-bit exponent, 27-bit significand¹
 - IBM: 32 bits, 1-bit sign, 7-bit exponent, 24-bit significand²
- Manufacturers implemented circuits with different rounding behavior, with/without infinity, and other inconsistencies
- Troublesome for reliability: code produced different results on different machines
- This was resolved with the adoption of the IEEE 754 Floating Point Standard which specifies
 - Bit layout of 32-bit float and 64-bit double
 - Rounding behavior, special values like Infinity

► Turing Award to William Kahan for his work on the standard

¹Floating Point Arithmetic ²IBM Hexadecimal Floats

IEEE 754 Format: The Standard for Floating Point

| float | double | Property |
|-------|--------|---|
| 32 | 64 | Total bits |
| 1 | 1 | Bits for sign (1 neg / 0 pos) |
| 8 | 11 | Bits for Exponent multiplier (power of 2) |
| 23 | 52 | Bits for Fractional part or mantissa |
| 7.22 | 15.95 | Decimal digits of accuracy ³ |

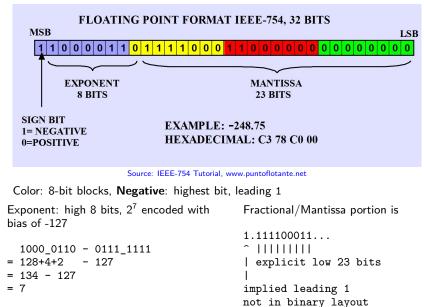
 Most commonly implemented format for floating point numbers in hardware to do arithmetic: processor has physical circuits to add/mult/etc. for this bit layout of floats

Numbers/Bit Patterns divided into three categories

| Category | Description | Exponent |
|--------------|---------------------------------------|-------------|
| Normalized | most common like 1.0 and -9.56e37 | mixed $0/1$ |
| Denormalized | very close to zero and 0.0 | all O's |
| Special | extreme/error values like Inf and NaN | all 1's |

³Wikipedia: IEEE 754

Example float Layout of -248.75: float_examples.c



Normalized Floating Point: General Case

- A "normalized" floating point number is in the standard range for float/double, bit layout follows previous slide
- ► Example: -248.75 = -1.111100011 * 2⁷

Exponent is in **Bias Form** (not Two's Complement)

- Unsigned positive integer minus constant bias number
- Consequence: exponent of 0 is not bitstring of 0's
- Consequence: tiny exponents like -125 close to bitstring of 0's; this makes resulting number close to 0
- 8-bit exponent 1000 0110 = 128+4+2 = 134 so exponent value is 134 - 127 = 7

Integer and Mantissa Parts

- The leading 1 before the binary point is **implied** so does not show up in the bit string
- Remaining fractional/mantissa portion shows up in the low-order bits

Fixed Bit Standards for Floating Point

IEEE Standard Layouts

| Kind | Sign | Exponent | | | Mantissa |
|--------|--------|-----------------|-------|------------------|----------------|
| | Bit | Bits | Bias | Exp Range | Bits |
| float | 31 (1) | 30-23 (8 bits) | -127 | -126 to +127 | 22-0 (23 bits) |
| double | 63 (1) | 62-52 (11 bits) | -1023 | -1022 to $+1023$ | 51-0 (52 bits) |

Standard allows hardware to be created that is as efficient as possible to do calculation on these numbers

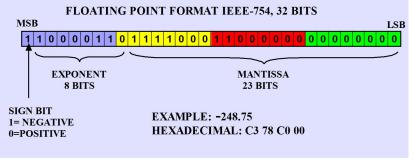
Consequences of Fixed Bits

- Since a fixed # of bit is used, some numbers cannot be exactly represented, happens in any numbering system:
- Base 10 and Base 2 cannot represent $\frac{1}{3}$ in finite digits

Try show_float.c to see this in action

Exercise: Quick Checks

- 1. What distinct parts are represented by bits in a floating point number (according to IEEE)
- 2. What is the "bias" of the exponent for 32-bit floats
- 3. Represent 7.125 in binary using "binary point" notation
- 4. Lay out 7.125 in IEEE-754 format
- 5. What does the number 1.0 look like as a float?



Source: IEEE-754 Tutorial, www.puntoflotante.net

The diagram above may help in recalling IEEE 754 layout

Special Cases: See float_examples.c

Special Values

- Infinity: exponent bits all 1, fraction all 0, sign bit indicates $+\infty$ or $-\infty$
- Infinity results from overflow/underflow or certain ops like float x = 1.0 / 0.0;
- #include <math.h> gets macro INFINITY and -INFINITY
- ▶ NaN: not a number, exponent bits all 1, fraction has some 1s
- Errors in floating point like 0.0 / 0.0

Denormalized values: Exponent bits all 0

- Fractional/Mantissa portion evaluates without implied leading one, still an unsigned integer though
- Exponent is Bias + 1: 2⁻¹²⁶ for float
- Result: very small numbers close to zero, smaller than any other representation, degrade uniformly to 0
- Zero: bit string of all 0s, optional leading 1 (negative zero);

Other Float Notes



Source: XKCD #217

Approximations and Roundings

Approximate ²/₃ with 4 digits, usually 0.6667 with standard rounding in base 10

- Similarly, some numbers cannot be exactly represented with fixed number of bits: ¹/₁₀ approximated
- IEEE 754 specifies various rounding modes to approximate numbers

Clever Engineering

- IEEE 754 allows floating point numbers to sort using signed integer sorting routines
- Bit patterns for float follows are ordered nearly the same as bit patterns for signed int
- Integer comparisons are usually fewer clock cycles than floating comparisons

Sidebar: The Weird and Wonderful Union

- Bitwise operations like & are not valid for float/double
- Can use pointers/casting to get around this OR...
- Use a union: somewhat unique construct to C
- Defined like a struct with several fields
- BUT fields occupy the same memory location (!?!)
- Allows one to treat a byte position as multiple different types, ex: int / float / char[]
- Memory size of the union is the max of its fields

```
// union.c
typedef union { // shared memory
float fl; // an float
int in; // a int
char ch[4]; // char array
} flint_t; // 4 bytes total
int main(){
flint_t flint;
flint.in = 0xC378C000;
```

```
}
```

| Ļ | Symbol | I | Mem | I | Val | 1 |
|----|-------------------|-----|-------|-----|------|----|
| 1. | | -+- | | -+- | | -1 |
| Т | flint.ch[3] | Ι | #1027 | Ι | 0xC3 | I |
| Т | flint.ch[2] | Ι | #1026 | Ι | 0x78 | I |
| Т | flint.ch[1] | Ι | #1025 | Ι | 0xC0 | I |
| L | flint.in/fl/ch[0] | Т | #1024 | Τ | 0x00 | I |
| L | i | T | #1020 | T | ? | I |

Floating Point Operation Efficiencies

- Floating Point Operations per Second, FLOPS is a major measure for numerical code/hardware efficiency
- Often used to benchmark and evaluate scientific computer resources, (e.g. top super computers in the world)
- Tricky to evaluate because of
 - A single FLOP (add/sub/mul/div) may take 3 clock cycles to finish: latency 3
 - Another FLOP can start before the first one finishes: pipelined
 - Enough FLOPs lined up can get average 1 FLOP per cycle
 - FP Instructions may automatically operate on multiple FPs stored in memory to feed pipeline: vectorized ops
 - Generally referred to as **superscalar**
 - Processors schedule things out of order too
- All of this makes micro-evaluation error-prone and pointless
- Run a real application like an N-body simulation and compute

 $\mathsf{FLOPS} = \frac{\mathsf{number of floating ops done}}{\mathsf{time taken in seconds}}$

| Rank | System | #Cores | Rmax (TFlop/s) | Rpeak (TFlop/s) | Power [*] (kW) |
|------|--|-----------|-------------------|--------------------|----------------------------|
| 1 | Frontier, <i>USA / Oak Ridge</i> Cray EX235a, AMD EPYC 2GHz (x86-64) | 8,730,112 | 1,102.00 | 1,685.65 | 21,100 |
| 2 | Fugaku, <i>Japan / Fujitsu</i> Fujitsu A64FX 2.2GHz (Arm) | 7,630,848 | 442,010.0 | 537,212.0 | 29,899 |
| 3 | LUMI <i>Finland / EuroHPC</i> Cray EX235a, AMD EPYC 2GHz (x86-64) | 1,110,144 | 151.90 | 214.35 | 2,942 |
| 4 | Summit <i>United States</i> IBM POWER9 22C 3.07GHz (Power) | 2,414,592 | 148,600.0 | 200,794.9 | 10,096 |
| 5 | Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power) | 1,572,480 | 94,640.0 | 125,712.0 | 7,438 |

https://www.top500.org/lists/top500/2022/06/

*: An average US Home uses 909 kWh of power per month

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| 3 | Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power) | 1,572,480 | 94,640.0 | 125,712.0 | 7,438 |
| 4 | Sunway TaihuLight <i>China</i> Sunway SW26010 (custom RISC) | 10,649,600 | 93,014.6 | 125,435.9 | 15,371 |
| 5 | Perlmutter, <i>United States</i> AMD EPYC 2.45GHz, Cray (x86-64) | 706,304 | 64,590.0 | 89,794.5 | 2,528 |

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| | | | Rmax | Rpeak | Power |
|------|--|------------|-----------|-----------|--------|
| Rank | System | #Cores | (TFlop/s) | (TFlop/s) | (kW) |
| 1 | Fugaku, <i>Japan / Fujitsu</i> Fujitsu A64FX 2.2GhZ (Arm) | 7,299,072 | 415,530.0 | 513,854.7 | 28,335 |
| 2 | Summit <i>United States</i> IBM POWER9 22C 3.07GHz (Power) | 2,397,824 | 143,500.0 | 200,794.9 | 10,096 |
| 3 | Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power) | 1,572,480 | 94,640.0 | 125,712.0 | 7,438 |
| 4 | Sunway TaihuLight <i>China</i> Sunway SW26010 (custom RISC) | 10,649,600 | 93,014.6 | 125,435.9 | 15,371 |
| 5 | Selene <i>USA, NVIDIA/AMD</i> AMD EPYC 7742 64C 2.25GHz (x86-64) | 555,520 | 63,460.0 | 79,215.0 | 2,646 |

https://www.top500.org/lists/top500/2020/06/

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| 5 | Tianhe-2A <i>China</i> Intel Xeon 2.2GHz (x86-64) | 4,981,760 | 61,444.5 | 100,678.7 | 18,482 |

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|------|---|------------|-----------|-----------|--------|
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| 1 | Summit <i>United States</i> IBM POWER9 22C 3.07GHz | 2,397,824 | 143,500.0 | 200,794.9 | 9,783 |
| 2 | Sierra <i>United States</i> IBM POWER9 22C 3.1GHz, | 1,572,480 | 94,640.0 | 125,712.0 | 7,438 |
| 3 | Sunway TaihuLight <i>China</i> Sunway MPP | 10,649,600 | 93,014.6 | 125,435.9 | 15,371 |
| 4 | Tianhe-2A <i>China</i> Xeon 2.2GHz | 4,981,760 | 61,444.5 | 100,678.7 | 18,482 |
| 5 | Frontera, <i>United States</i> Dell 6420, Xeons 2.7GHz | 448,448 | 23,516.4 | 38,745.9 | ?? |

https://www.top500.org/list/2019/11/

| | | | Rmax | Rpeak | Power |
|------|---|------------|-----------|-----------|--------|
| Rank | System | #Cores | (TFlop/s) | (TFlop/s) | (kW) |
| 1 | Summit <i>United States</i> IBM POWER9 22C 3.07GHz | 2,397,824 | 143,500.0 | 200,794.9 | 9,783 |
| 2 | Sierra <i>United States</i> IBM POWER9 22C 3.1GHz, | 1,572,480 | 94,640.0 | 125,712.0 | 7,438 |
| 3 | Sunway TaihuLight <i>China</i> Sunway MPP | 10,649,600 | 93,014.6 | 125,435.9 | 15,371 |
| 4 | Tianhe-2A <i>China</i> TH-IVB-FEP Cluster | 4,981,760 | 61,444.5 | 100,678.7 | 18,482 |
| 5 | Piz Daint <i>Switzerland</i> Cray XC50, Xeon E5-2690v3 | 387,872 | 21,230.0 | 27,154.3 | 2,384 |

https://www.top500.org/list/2018/11/

| Rank | System | #Cores | Rmax (TFlop/s) | Rpeak (TFlop/s) | Power (kW) |
|------|--|------------|-------------------|--------------------|---------------|
| 1 | Sunway TaihuLight <i>China</i> Sunway MPP | 10,649,600 | 93,014.6 | 125,435.9 | 15,371 |
| 2 | Tianhe-2 (MilkyWay-2) <i>China</i> TH-IVB-FEP Cluster | 3,120,000 | 33,862.7 | 54,902.4 | 17,808 |
| 3 | Piz Daint <i>Switzerland</i> Cray XC50 | 361,760 | 19,590.0 | 25,326.3 | 2,272 |
| 4 | Gyoukou <i>Japan</i> ZettaScaler-2.2 HPC system | 19,860,000 | 19,135.8 | 28,192.0 | 1,350 |
| 5 | Titan <i>USA</i> Cray XK7 | 560,640 | 17,590.0 | 27,112.5 | 8,209 |

https://www.top500.org/lists/2017/11/